

# Algorithmic Game Theory

Winter Term 2022/2023

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## Assignment 8

Issued: Dec 13, 2022  
Due: Dec 20, 2022, 10:00h

This and all upcoming assignments are only relevant for the second part of the course.

### Exercise 8.1.

(2 + 2 Points)

There is an auction with one good and  $n > 2$  bidders with private valuations for the good. We assume that their bids are pairwise distinct. The good is assigned to the bidder with the highest bid, and the price is  $p$ . The other bidders neither receive nor pay anything.

State all values of  $n$  for which the described mechanism is incentive compatible, when  $p$  is

- the arithmetic mean of all bids,
- the (lower) median bid (always take the  $\lfloor \frac{n+1}{2} \rfloor$ -th smallest bid).

Prove your answers.

### Exercise 8.2.

(4 Points)

We consider a combinatorial auction: There is a set  $\mathcal{G}$  of goods that are indivisible. Bidders have a private valuation for each subset of  $\mathcal{G}$ . Every outcome allocates exactly one subset of  $\mathcal{G}$  to every bidder such that the intersection of each pair of allocated subsets is empty.

Let  $\mathcal{G} = \{A, B, C\}$  and  $\mathcal{N} = \{1, 2, 3, 4\}$ . The private valuations  $v_i(G)$  for all bidders  $i \in \mathcal{N}$  and subsets  $G \subseteq \mathcal{G}$  are given by the following table:

	$\emptyset$	$\{A\}$	$\{B\}$	$\{C\}$	$\{A, B\}$	$\{A, C\}$	$\{B, C\}$	$\{A, B, C\}$
$i = 1$	0	2	3	2	2	4	6	6
$i = 2$	0	0	0	3	0	6	5	9
$i = 3$	0	1	5	1	4	1	1	2
$i = 4$	0	1	2	3	3	5	3	5

Apply the VCG mechanism on this instance: Explain which allocation is chosen by VCG, and why. Then calculate the payments for all bidders (with explanation). You may assume true bidding.

**Exercise 8.3.**

(3 + 3 + 2 Points)

For a given undirected graph  $G = (V, E)$  with edge capacities  $c_e \in \mathbb{N}$  for all  $e \in E$ , we want to auction matchings. Each bidder  $i \in \mathcal{N}$  has a desired matching  $M_i \subseteq E$ , a demand  $d_i \in \mathbb{N}$  and a private valuation  $v_i \in \mathbb{R}_+$  when they are selected. Otherwise, the valuation is 0.

After seeing all bids, the auctioneer selects a subset  $\mathcal{I} \subseteq \mathcal{N}$  of bidders and decides on the prices for all bidders. Let  $M$  be the union of matchings from selected bidders, i.e.  $M = \bigcup_{i \in \mathcal{I}} M_i$ . The selection of bidders  $\mathcal{I}$  needs to be feasible in the following sense:

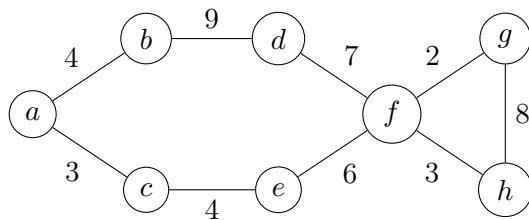
- $M$  has to be a matching itself.
- For each  $e \in M$ , the sum of demands of all  $i \in \mathcal{I}$  with  $e \in M_i$  does not exceed  $c_e$ .

A selection rule for this auction when receiving bids  $b_1, \dots, b_n$  is given as follows: We start with  $\mathcal{I} = \emptyset$  and remove bidders with  $M_i = \emptyset$  or  $b_i = 0$  beforehand. Then:

- Sort bidders by their reported value per demand per edge, i.e. by  $\frac{b_i}{d_i \cdot |M_i|}$ . Break ties arbitrarily.
- Iterate over the ordered bidders and greedily add bidder  $i$  into  $\mathcal{I}$ , unless the selection becomes infeasible due to  $i$ .

a) Prove that the given selection rule can be used by an incentive compatible mechanism. Construct such mechanism, i.e. state a suitable payment formula.

b) Consider the following instance with 4 bidders:



$i$	$M_i$	$d_i$	$v_i$
1	$\{\{a, b\}, \{c, e\}, \{g, h\}\}$	3	27
2	$\{\{d, f\}, \{g, h\}\}$	5	30
3	$\{\{a, c\}, \{b, d\}, \{e, f\}, \{g, h\}\}$	3	24
4	$\{\{c, e\}, \{d, f\}\}$	2	10

Calculate the selection and the prices for all bidders when the mechanism from a) is applied.

*Hint:* Pictures might be helpful when explaining your calculation, which is required!

c) Add new bidders to the instance from b) such that the following requirements are met:

- There are at most two additional bidders.
- The selection retrieved by the mechanism from a) is not socially optimal.

Prove the correctness of your modification.